

# GENERALIZED CRUDE BRAUER INEQUALITY ON ADDITION CHAINS

THEOPHILUS AGAMA

ABSTRACT. We extend the inequality due to Alfred Brauer on standard addition chains to a sequence of additions leading to a finite number where at most at most  $d \geq 2$  previous terms can be added to generate each term in the sequence.

## 1. INTRODUCTION

An addition chain of length  $h$  leading to  $n$ , first introduced by the German mathematician Arnold Scholz in 1937 [1], is a sequence of numbers  $s_0 = 1, s_1 = 2, \dots, s_h = n$  where  $s_i = s_k + s_s$  for  $i > k \geq s \geq 0$ . The number of terms (excluding the first term) in an addition chain leading to  $n$  is the length of the chain. We call an addition chain leading to  $n$  with a minimal length an *optimal* addition chain leading to  $n$ . In standard practice, we denote by  $\ell(n)$  the length of an optimal addition chain that leads to  $n$ .

**Example 1.1.** Consider the addition chain  $\mathcal{C}_{16} : s_0 = 1, s_1 = 2, s_2 = 4, s_3 = 8, s_4 = 9, s_5 = 16$  with sums

$$\begin{aligned} s_1 &= 2 = 1 + 1 \\ s_2 &= 4 = 2 + 2 \\ s_3 &= 8 = 4 + 4 \\ s_4 &= 9 = 8 + 1 \\ s_5 &= 16 = 8 + 8. \end{aligned}$$

In [1], Arnold Scholz asked if the crude inequality  $\ell(n) \leq 2 \frac{\log n}{\log 2}$  could be improved. In the paper [2], Alfred Brauer proved the class of inequality:

---

*Date:* December 21, 2025.

*2010 Mathematics Subject Classification.* Primary 11P32, 11A41; Secondary 11B13, 11H99.

*Key words and phrases.* Generalized, Brauer.

**Theorem 1.2.**

$$\ell(n) \leq \min_{1 \leq r \leq m} \left[ \left(1 + \frac{1}{r}\right) \frac{\log n}{\log 2} + 2^r - 2 \right]$$

for  $2^m \leq n < 2^{m+1}$ .

Brauer obtained Theorem 1.2 by showing that there exists an addition chain leading to  $n$  for  $2^{rs} \leq n < 2^{r(s+1)}$ , starting with  $1, 2, \dots, 2^r - 1$  and having length at most  $(1+r)s + 2^r - 2$ . Fixing  $r$ , he argued basically by the method of contradiction with  $s$ . Using the inequality  $s \leq \frac{\log n}{r \log 2}$ , which can easily be deduced from the constraint on  $n$ , one obtains Theorem 1.2. By choosing  $r := \lfloor \log \log n \rfloor + 1$  in Theorem 1.2 - which comes from optimizing the quantity on the right-hand side - he obtained the upper bound as a consequence

*Proof of Theorem 1.2.* Theorem 1.2 is obtained in the spirit of showing that for any positive integer  $r$  and any non-negative integer  $s$  there exists an addition chain leading to  $2^{rs} \leq n < 2^{r(s+1)}$  that starts with the first  $2^r - 1$  consecutive positive integers  $1, 2, \dots, 2^r - 1$  and has length

$$\ell(n) \leq (1+r)s + 2^r - 2.$$

In the case  $r = 1$ , we obtain as a consequence

$$\ell(n) \leq 2s$$

for  $2^s \leq n < 2^{s+1}$ . The claim holds for this choice of  $r$  since this is the right-hand side inequality of Arnold Scholz established by Alfred Brauer.

In the case  $s = 0$ , we obtain the statement: There exists an addition chain that leads to  $n < 2^r$  starting with the first  $2^r - 1$  consecutive integers and has length  $\ell(n) \leq 2^r - 2$ . This is obviously true since with the target  $n = 2^r - 1$  the first  $2^r - 1$  consecutive integers is an addition chain leading to  $2^r - 1$  with length  $2^r - 2$ . This serves as an upper bound for the optimal length of all targets  $n < 2^r$ .

Fix  $r > 1$  and suppose that there exists some  $s$  for which there are no addition chains leading to  $n \in [2^{rs}, 2^{r(s+1)})$  satisfying the length requirement  $\ell(n) \leq (1+r)s + 2^r - 2$  and starts with the first  $2^r - 1$  consecutive integers. Let  $s_o$  be the smallest among those  $s$  that satisfy this observation. This implies that there are no addition chains leading to  $n$  with  $2^{rs_o} \leq n < 2^{r(s_o+1)}$  satisfying  $\ell(n) \leq (1+r)s_o + 2^r - 2$  and containing the first  $2^r - 1$  consecutive positive integers. Now divide  $n$  by  $2^r$ , then we can write  $n = a2^r + b$  for  $2^{r(s_o-1)} \leq a < 2^{rs_o}$  and  $b < 2^r$ . It follows that there is an addition chain

$$a_o = 1, a_1 = 2, \dots, a_l = a$$

that starts with the first  $2^r - 1$  consecutive positive integers and has length  $\ell(a) \leq (1+r)(s_o - 1) + 2^r - 2$ .

We now construct an addition chain leading to  $n$  using the addition chain leading to  $a$  as follows

$$1 = a_o, 2 = a_1, \dots, a_l = a, 2a, \dots, 2^r a, 2^r a + b = n$$

since  $b \leq 2^r - 1$  and must be a term in the chain leading to  $a$ . The length of this addition chain is at most

$$(1+r)(s_o - 1) + (r+1) + 2^r - 2 = (1+r)s_o + 2^r - 2$$

which violates the earlier assumption that  $s_o$  is the smallest positive integer satisfying the observation. Therefore, the claim holds for all nonnegative integers  $s$  for a fixed  $r$ . Using the inequality  $s \leq \frac{\log n}{r \log 2}$ , we deduce

$$\ell(n) \leq \left(1 + \frac{1}{r}\right) \frac{\log n}{\log 2} + 2^r - 2$$

for all  $r > 1$ . The inequality holds for the minimum over all  $r > 1$ . Optimizing

$$\left(1 + \frac{1}{r}\right) \frac{\log n}{\log 2} + 2^r - 2$$

we obtain an ideal choice for  $r := \lfloor \log \log n \rfloor + 1$  which yields the following inequality.

**Corollary 1.3.**

$$\ell(n) < \frac{\log n}{\log 2} \left\{ 1 + \frac{1}{\log \log n} + \frac{2 \log 2}{(\log n)^{1-\log 2}} \right\}$$

for  $n \geq n_o$  for some fixed  $n_o > 0$

a result which significantly improves the upper bound  $\ell(n) \leq 2 \frac{\log n}{\log 2}$  for all sufficiently large  $n$ .

2. A CRUDE GENERALIZATION

**Definition 2.1.** Let  $n \geq 3$  and  $d \geq 2$  be positive integers. We say that the sequence of positive integers

$$s_0 = 1 < s_1 < \dots < s_h = n$$

is an addition chain with fixed degree  $d \geq 2$  leading to  $n$  of length  $h$  if for each  $1 \leq i \leq h$  the representation

$$s_i = \sum_{\substack{i_j \in [0, i-1] \cap \mathbb{Z} \\ 1 \leq j \leq d}} s_{i_j}, \quad (s_{i_j} < s_i)$$

holds.

In other words, each term in an addition chain with a fixed degree  $d \geq 2$  is the sum of at most  $d$  previous terms in the chain, with repetition allowed. We call the shortest addition chain with fixed degree  $d \geq 2$  leading to a target an *optimal* addition chain with fixed degree  $d$ . We denote the length of an optimal addition chain with fixed degree  $d \geq 2$  leading to a target  $n$  with  $\ell^d(n)$ . The special case where the fixed degree  $d = 2$  recovers the well-known concept of an addition chain.

**Example 2.2.** Choose the target  $n = 21$  and fix the degree  $d = 3$ . The sequence

$$s_0 = 1, s_1 = 2, s_2 = 4, s_3 = 8, s_4 = 16, s_5 = 21$$

is an addition chain with fixed degree  $d = 3$ , because  $s_2 = 2s_1, s_3 = 2s_2, s_4 = 2s_3, s_5 = s_4 + s_2 + s_0$ . However, it is not of minimal length. An example chain of fixed degree 3 and of minimal length is

$$s_0 = 1, s_1 = 3, s_2 = 9, s_3 = 21.$$

**Theorem 2.3** (Generalized Brauer-type bound). *Let  $d \geq 2$  be fixed and let  $n$  be a positive integer satisfying  $d^m \leq n < d^{m+1}$ . Then*

$$\ell^d(n) \leq \min_{1 \leq r \leq v \leq m} \left[ \left(1 + \frac{1}{r}\right) \frac{\log n}{\log d} + d^v - 2 \right].$$

*Proof.* Let  $d \geq 2$  be fixed and let  $n$  be a positive integer such that  $d^m \leq n < d^{m+1}$ . We choose  $r \geq 1$  to be a positive integer and set

$$k := \left\lceil \frac{m}{r} \right\rceil$$

and

$$v := \max_{0 \leq i \leq k-1} (r(i+1) - i).$$

Let  $u_1, u_2, \dots, u_k$  be a sequence of positive integers such that  $0 \leq u_j < d^v$  for  $1 \leq j \leq k$  and such that

$$n = \sum_{j=0}^{k-1} u_{k-j} d^{r(k-(j+1))}$$

where

$$u_{k-j} < d^{r(j+1)-j} \leq d^v.$$

This representation exists and is unique. We check this by considering the base- $d$  expansion of  $n$  as

$$n = a_0 + a_1 d + \dots + a_m d^m$$

where  $0 \leq a_i < d$ . Because  $(k-1)r \leq m < kr$ , the representation can be rewritten in the form

$$n = b_0 + b_1 d^{kr-(kr-1)} + \dots + b_p d^{kr-1}$$

with  $0 \leq b_i < d$  for  $0 \leq i \leq p$ . We can therefore write

$$\begin{aligned} b_{p-j} d^{kr-(j+1)} &= b_{p-j} d^{r(k-(j+1))} d^{r(j+1)-(j+1)} \\ &= d^{r(k-(j+1))} u_{k-j} \end{aligned}$$

where

$$u_{k-j} := b_{p-j} d^{r(j+1)-(j+1)} < d^{r(j+1)-j}.$$

Therefore, the representation

$$n = \sum_{j=0}^{k-1} u_{k-j} d^{r(k-(j+1))}$$

with

$$u_{k-j} := b_{p-j} d^{r(j+1)-(j+1)} < d^{r(j+1)-j}$$

exists for each  $n$  such that  $d^m \leq n < d^{m+1}$  and is unique. Now, we construct an addition chain with fixed degree  $d \geq 2$  starting with the sequence of consecutive integers

$$s_0 = 1, s_1 = 2, \dots, s_{d^v-2} = d^v - 1$$

and continue the degree  $d$  chain as

$$\begin{aligned} s_{d^v-1} &= u_k, \\ s_{d^v-1+r} &= d^r u_k, \\ s_{d^v+r} &= d^r u_k + u_{k-1}, \\ s_{d^v+2r} &= d^{2r} u_k + d^r u_{k-1}, \\ s_{d^v+2r+1} &= d^{2r} u_k + d^r u_{k-1} + u_{k-2}. \end{aligned}$$

Continue the degree  $d \geq 2$  addition chain in a similar fashion until we reach

$$s_{d^v+(k-1)r+k-2} = \sum_{j=0}^{k-1} u_{k-j} d^{r(k-(j+1))} = n.$$

It follows that  $\ell^d(n) \leq d^v - 2 + (k-1)r + (k-1) \leq d^v - 2 + m + \frac{m}{r}$ . Using the fact  $m \leq \frac{\log n}{\log d}$  and the inequality holds for all  $r \geq 1$ , the claim follows immediately.  $\square$

We can specialize Theorem 2.3 by taking  $v = r$  and setting

$$F(r) := \left[ \left( 1 + \frac{1}{r} \right) \frac{\log n}{\log d} + d^r - 2 \right]$$

as a continuous function of  $r$ . We minimize  $F(r)$  by setting  $\frac{dF}{dr} = 0$ . A quick elementary calculation yields

$$\frac{dF}{dr} = -\frac{1}{r^2} \frac{\log n}{\log d} + d^r \log d = 0$$

which gives the equation

$$r \log d + 2 \log \log d = \log \log n - 2 \log r.$$

The integer choice of  $r$  which gives the local minimum for  $F(r)$  can be approximated by

$$r := \left\lfloor \frac{\log \log n}{\log d} \right\rfloor.$$

By plugging this choice of  $r$  into Theorem 2.3, we deduce the general upper bound

**Corollary 2.4.** *Let  $d \geq 2$  be fixed and denote  $\{\cdot\}$  as the fractional part of  $\cdot$ . Then*

$$\ell^d(n) < \frac{\log n}{\log d} \left( 1 + \frac{\log d}{(\log 2)(\log \log n)} + \frac{\log d}{d^{\{\frac{\log n}{\log d}\}}} \right)$$

for all  $n \geq n_0$  for some fixed  $n_0 > 0$ .

*Remark 2.5* (Improvement communicated by N. Doyon). Immediately after obtaining this result, the author received a private communication from his research director, N. Doyon, containing a substantially stronger upper bound than the one stated in Theorem 2.3. The bound communicated to the author dramatically improves the asymptotic upper bound for  $\ell^d(n)$ . The sharper inequality communicated is a manuscript in preparation for a journal submission and therefore kept confidential.

---

## REFERENCES

1. A. Scholz, *253, Jber. Deutsch. Math. Verein.* II, vol. 47, 1937, 41–42.
2. A. Brauer, *On addition chains*, Bulletin of the American mathematical Society, vol. 45:10, 1939, 736–739.